

Prototypical Problem

$$\begin{aligned} \max: \quad & Z = C_1 x_1 + \dots + C_n x_n \\ \text{Sub. to:} \quad & \begin{cases} a_{11}x_1 + \dots + a_{1n}x_n \leq b_1 \\ \vdots \\ a_{m1}x_1 + \dots + a_{mn}x_n \leq b_m \\ x_i \geq 0, \quad 1 \leq i \leq n \end{cases} \quad (*) \end{aligned}$$

Equivalent Problem ^①

$$\begin{aligned} \max: \quad & Z = C_1 x_1 + \dots + C_n x_n \\ \text{Sub. to:} \quad & \begin{cases} a_{11}x_1 + \dots + a_{1n}x_n + x_{n+1} = b_1 \\ \vdots \\ a_{m1}x_1 + \dots + a_{mn}x_n + x_{n+m} = b_m \\ x_i \geq 0, \quad x_{n+j} \geq 0, \quad 1 \leq i \leq n, \quad 1 \leq j \leq m \end{cases} \quad (*) \end{aligned}$$

Properties:

- Constraint (*) defines an n -d simplex in $(n+m)$ -d space $(x_1, \dots, x_{n+m})$.
- Max. of Z occurs at a vertex, or corner point of the simplex.
- The simplex is bounded by $(n+m-1)$ -d hyperplanes, and each hyperplane is defined by the equation $x_i = 0$ for some $i = 1, 2, \dots, n, n+1, \dots, n+m$.
- Each vertex or corner point of the n -d simplex (feasible region) is the intersection of n hyperplanes, i.e., $x_{i_1} = x_{i_2} = \dots = x_{i_n} = 0$ for $1 \leq i_k \leq n+m$. Its complete characterization, $(x_1, x_2, \dots, x_n)$, requires finding the remaining values, ~~for~~ x_j with $j \neq i_k$, from the system of ~~equations~~ the m constraint equations (*).
- ~~Therefore~~ The method of Gaussian elimination is used to solve the m constraint equations (*) for x_j , $j \neq i_k$, with $x_{i_1} = x_{i_2} = \dots = x_{i_n} = 0$.
- The simplex method is an iterative method ~~for~~ by which at each iteration it is to determine which corner point is to determine, or which ~~of~~ n variables should be assigned to zero (or the status of nonbasic variables) and then solve for the remaining m variables from the m many constraint equations.

Simplex Method (algebraic Form and Tabular Form) @ Standard (or canonical) Form.

$$\max: Z - c_1 x_1 - c_2 x_2 - \dots - c_n x_n = 0$$

$$\begin{cases} a_{11}x_1 + \dots + a_{1n}x_n + x_{n+1} = b_1 \\ \vdots \\ a_{m1}x_1 + \dots + a_{mn}x_n + \dots + x_{n+m} = b_m \end{cases}$$

Initial Corner Point Feasible: $x_1 = x_2 = \dots = x_n = 0$, $x_{n+1} = b_1, \dots, x_{n+m} = b_m$
 i.e.: $(0, \dots, 0, b_1, \dots, b_m)$ w/ Z -value: $Z = 0$

Use Gaussian Elimination to derive the k th canonical form:

$$\max: Z + c_{1,k}x_1 + c_{2,k}x_2 + \dots + c_{n+m,k}x_{n+m} = b_{0,k}$$

$$a_{11,k}x_1 + a_{12,k}x_2 + \dots + a_{1n,k}x_n + \dots + x_{n+1} = b_{1,k}$$

$\vdots$

$$a_{m1,k}x_1 + a_{m2,k}x_2 + \dots + a_{mn,k}x_n + \dots + x_{n+m} = b_{m,k}$$

w/ n hyperplanes (or nonbasic variables): $x_{i_1} = x_{i_2} = \dots = x_{i_n} = 0$

The form is canonical if (through by Gaussian Elimination) the following conditions are satisfied:

(1) All the coefficients of the basic variables in the ~~the Z -equation~~ the Z -equation are zero: $c_{j,k} = 0$ for $j \neq i_l, 1 \leq l \leq n$.

(2) For each basic variable $x_j, 1 \leq j \leq m, j \neq i_l, 1 \leq l \leq n$, there is one and only one constraint equation, $j_r, 1 \leq r \leq m$, so that all the coefficients of the basic variables ~~are zero except for the coefficient of x_j which is 1, i.e.~~ are zero except for the coefficient of x_j which is 1, i.e.

$$a_{j_r j,k} = 1, \text{ and } a_{j_r p,k} = 0, p \neq j, p \neq i_l, 1 \leq l \leq n.$$

(3) ~~As a result, the~~ All right hand values are ^{non-negative} ~~possible~~, $b_{i,k} \geq 0, 0 \leq i \leq m$.

(4) As a result, the corner point feasible: $x_{i_1} = x_{i_2} = \dots = x_{i_n} = 0$, $x_j = b_{j_r,k} \geq 0, j \neq i_l$, with Z -value: $Z = b_{0,k}$.

Criteria to stop the iterative process of simplex method. ⁽³⁾

- At the k th iterative canonical form

$$Z + c_{1,k}x_1 + c_{2,k}x_2 + \dots + c_{n+m,k}x_{n+m} = b_{0,k}$$

with nonbasic variables $x_{i_1} = x_{i_2} = \dots = x_{i_n} = 0$, and

basic variables $x_j > 0$ ~~with $j \neq i_l, 1 \leq l \leq n$~~ , for which $c_{j,k} = 0$ with $j \neq i_l, 1 \leq l \leq n$, if all the coefficients, $c_{j_l,k}$, for the nonbasic variables are nonnegative, then the k th Z -value, $b_{0,k}$, is the maximum because it cannot be improved further by increasing any x_{i_l} from 0.

So in this case, optimal solution is found, and stop the method.

- If there is ~~at least one coefficient~~ a negative coefficient, say $c_{i_l,k} < 0$, for at least one nonbasic variable x_{i_l} , the simplex method need continue for one more step.
- Here is the algorithm to determine the new set of nonbasic variables (i.e. new set of n hyperplanes for the new corner point which is adjacent to the old one).
 - let i_l be the index so that the coefficient $c_{i_l,k}$ ^{(i.e. $(k+1)$ st)} is the most negative so that the new set of nonbasic variables must NOT include x_{i_l} .
 - To determine the new set of nonbasic variables, we must find a variable from the old ^(i.e. k th) basic variables $x_j, j \neq i_0, 1 \leq j \leq n$, to take the place of x_{i_l} . This x_j is determined by the Ratio Test. Recall from the k th canonical form that the $x_j, j \neq i_0$, is solved from the j_r th equation of the constraints. The new nonbasic variable is this x_j so that $b_{j_r,k} / a_{j_r,k}$ is the smallest for $1 \leq r \leq m$. The new basic variable x_{i_l} now corresponds to the j_r th equation in the $(k+1)$ st iteration of the Gaussian elimination.