

1. Fix $m > 0$. The q -ary Hamming code C has parity-check matrix H , where H is an $m \times n$ matrix whose columns are a nonzero vector in each of the 1-dimensional vector spaces in \mathbb{F}_q^m (that is, no column is a multiple of another column). Determine the parameters n, k, d such that C is an $[n, k, d]$ code. Show that C is perfect.
2. Let C be an $[n, k, 3]$ GRS code with code locators $\alpha_1, \dots, \alpha_n \in \mathbb{F}_q$ and where the column multipliers are all 1. Suppose that codeword c is transmitted, y is received, and the error $e = y - c$ contains one error in position j . Let $(S_0, S_1) = Hy^T$ be the syndrome of y with respect to the canonical parity-check matrix H of C . Show that $\alpha_j = S_1/S_0$ and that the error value e_j is equal to S_0 .
3. Show that the minimum distance of a perfect code must be odd.
4. A burst of length ℓ is the event of having errors in a codeword such that the locations i and j of the first and last errors, respectively, satisfy $j - i = \ell - 1$. Let C be a linear $[n, k]$ code over \mathbb{F}_q and suppose that there exists a decoder for C that corrects every burst of length t or less.
 - (a) Show that in every nonzero codeword c in C , the locations i and j of the first and last nonzero entries in c must satisfy $j - i \geq 2t$.
 - (b) *The Reiger bound*: a Singleton-like bound for burst correcting codes. Show that $n - k \geq 2t$.
 - (c) A sphere-packing-like bound for burst-correcting codes. Show that

$$q^{n-k} \geq 1 + n(q-1) + (q-1)^2 \sum_{i=0}^{t-2} (n-i-1)q^i.$$

5. Let G be a $k \times n$ generator matrix of a linear code C over \mathbb{F}_q , where $k > 0$. Show that the minimum distance of C is the largest integer d such that every $k \times (n - d + 1)$ sub-matrix of G has rank k . (*Hint*: Show that if J is the set of indices of the zero entries in a nonzero codeword, then the columns in G that are indexed by J form a $k \times |J|$ submatrix whose rows are linearly dependent. Conversely, show that if the columns indexed by a set J form a submatrix whose rows are linearly dependent, then there is a nonzero codeword in which the entries that are indexed by J equal zero.)
6. Let C be a linear $[n, k, d]$ code over \mathbb{F}_q .
 - (a) Show that C is MDS if and only if every set of k columns in its generator matrix is linearly independent. (*Hint*: Use the previous problem.)
 - (b) Assume that $k < n$. Show that C is MDS if and only if its dual code is.